

# Cost-Aware Optimal Filter Assignment Policy Against Distributed Denial-of-Service Attack

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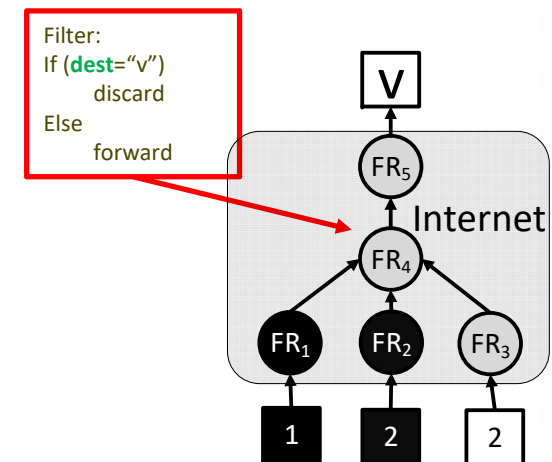
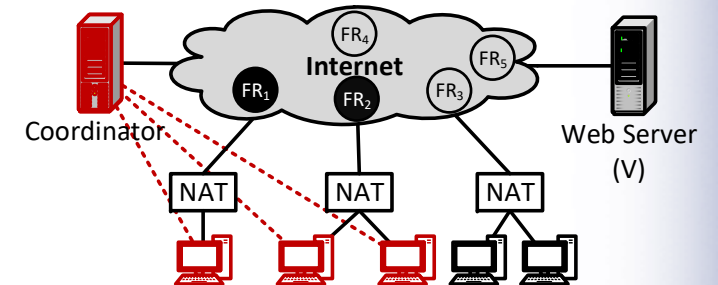
## Outline

- Introduction to DDoS attack and filter router
- Previous works
- Four phase DDoS protection system model
- Problem: Minimizing blocked legitimate user
- Dynamic programming solution
- Simulation results
- Q & A



# DDoS & Four-phase Protection System

- DDoS
  - Attacker keeps the victim busy.
  - Millions of requests are fired by bots.
  - Bots are controlled by a master.
  
- Background
  - Filter router
    - Does packet marking.
    - Apply filter and block traffic according to filter.
  - Filter
    - Simple packet blocking rule.
    - Source-based, **destination based**.



# Previous work

Systems	Limitations
<p><b>Probabilistic Filter Scheduling (packet marking)</b></p> <p>PFS: Probabilistic filter scheduling against distributed denial-of-service attacks (D. Seo et al. in IEEE 36th Conf. Local Comput. Netw, Oct. 2011)</p>	<ul style="list-style-type: none"> <li>• Does not consider limited budget on filters.</li> <li>• Filter propagation takes some time.</li> <li>• Hard to send huge number of filters.</li> </ul>
<p><b>Filter Scheduling (block all attack traffic)</b></p> <p>Filter Assignment Policy Against Distributed Denial-of-Service Attack (R. Biswas et al. in IEEE ICPADS, Dec. 2018, )</p>	<ul style="list-style-type: none"> <li>• Cannot assign filter optimally.</li> <li>• Blocking of all attack traffic increase blocking of legitimate users.</li> </ul>





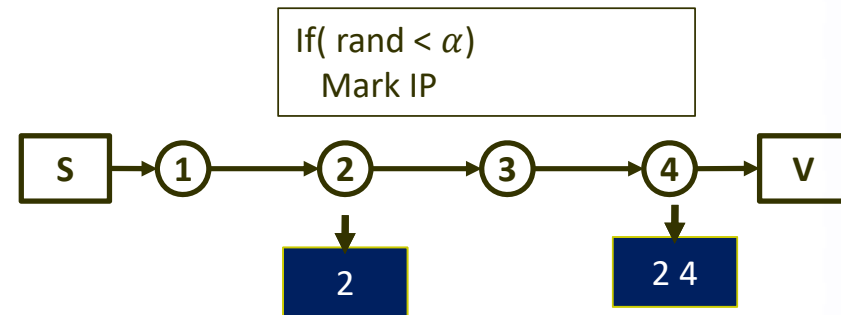
## A Four-phase Protection Process

- Phase I: **Packet marking** by Filter Router.
- Phase II: Traffic **topology** and **filter construction**.
- Phase III: **Assign filters** to filter router.
- Phase IV: **Evict unused filter** from filter router.

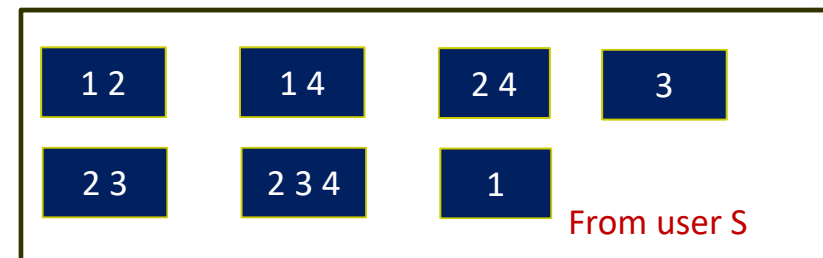


# Phase I: Packet Marking by FR

- Filter router (FR) probabilistically appends its own IP address to the packet.
- $\alpha$  = marking probability

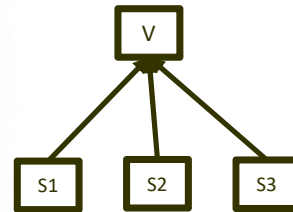


Example received packets,  $\alpha = 0.5$



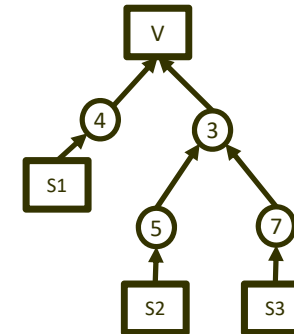
## Phase II: Topology Construction

- S1
- S2
- S3



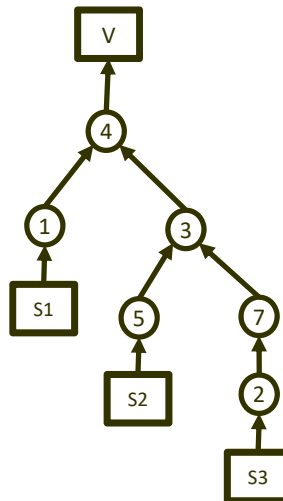
Without any marking

- S1 4
- S2 5 3
- S3 7 3



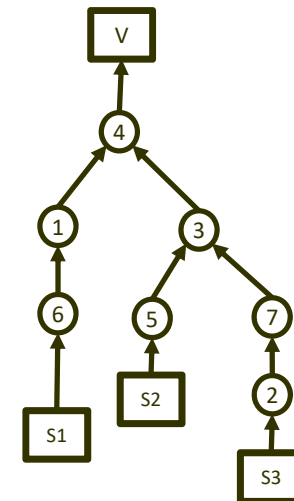
After few marking received

- S1 1 4
- S2 3 4
- S3 2 7



After few more marking received

- S1 6 4
- S2 4
- S3 3



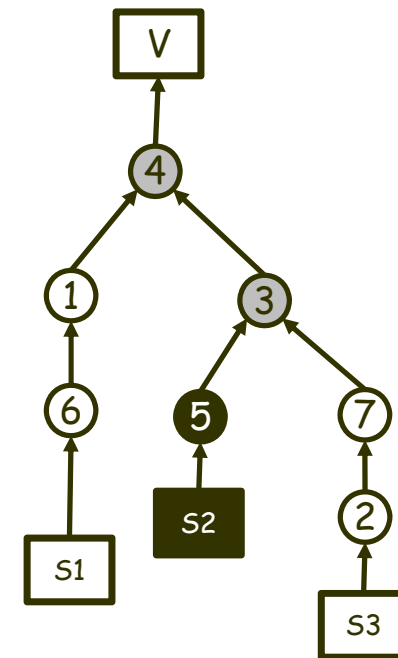
After some more marking received



## Identifying Attackers' IP

- Victim can identify attacker.
  - Statistical approaches, packet arrival time, entropy, etc.

- Black: only attacker traffic
- White: only legitimate traffic
- Gray: mixed traffic



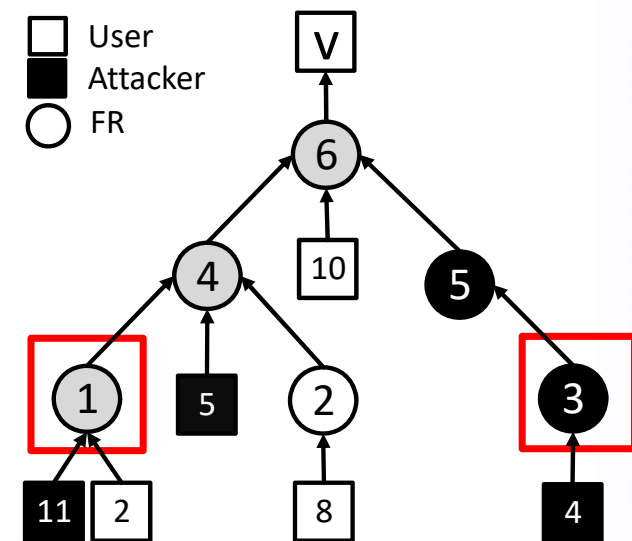
Sending filters to large # of FR takes long time. The ISP of FR may charge money.





## Problem: Minimizing Blocked Legit Users

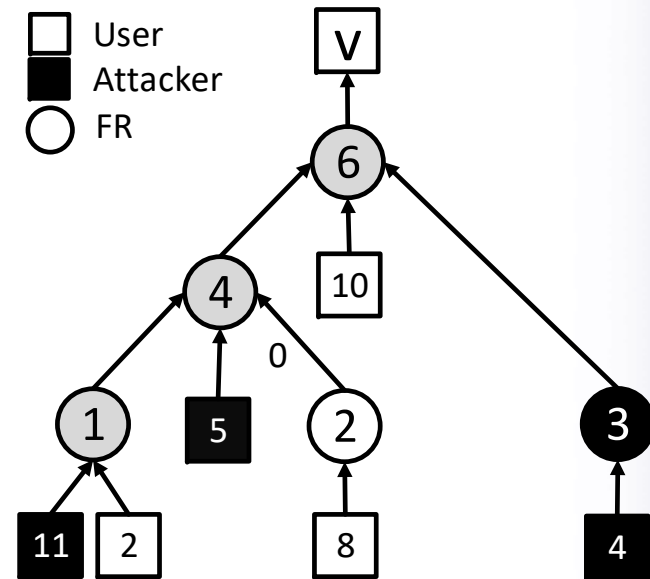
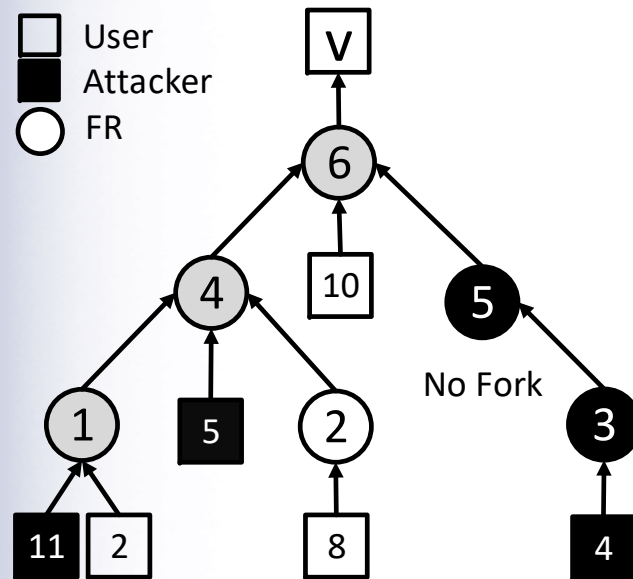
- Given topology, select  $K$  filters so that  $C$  is minimum.
- Cost model
  - $C$  = Number of blocked legit user
- Constraint
  - Yielded traffic  $\leq$  bandwidth of victim
- Best assignment for  $k=2$  is  $\{1,3\}$ 
  - $C = 2$ , Yielded traffic= 23



Bandwidth = 25



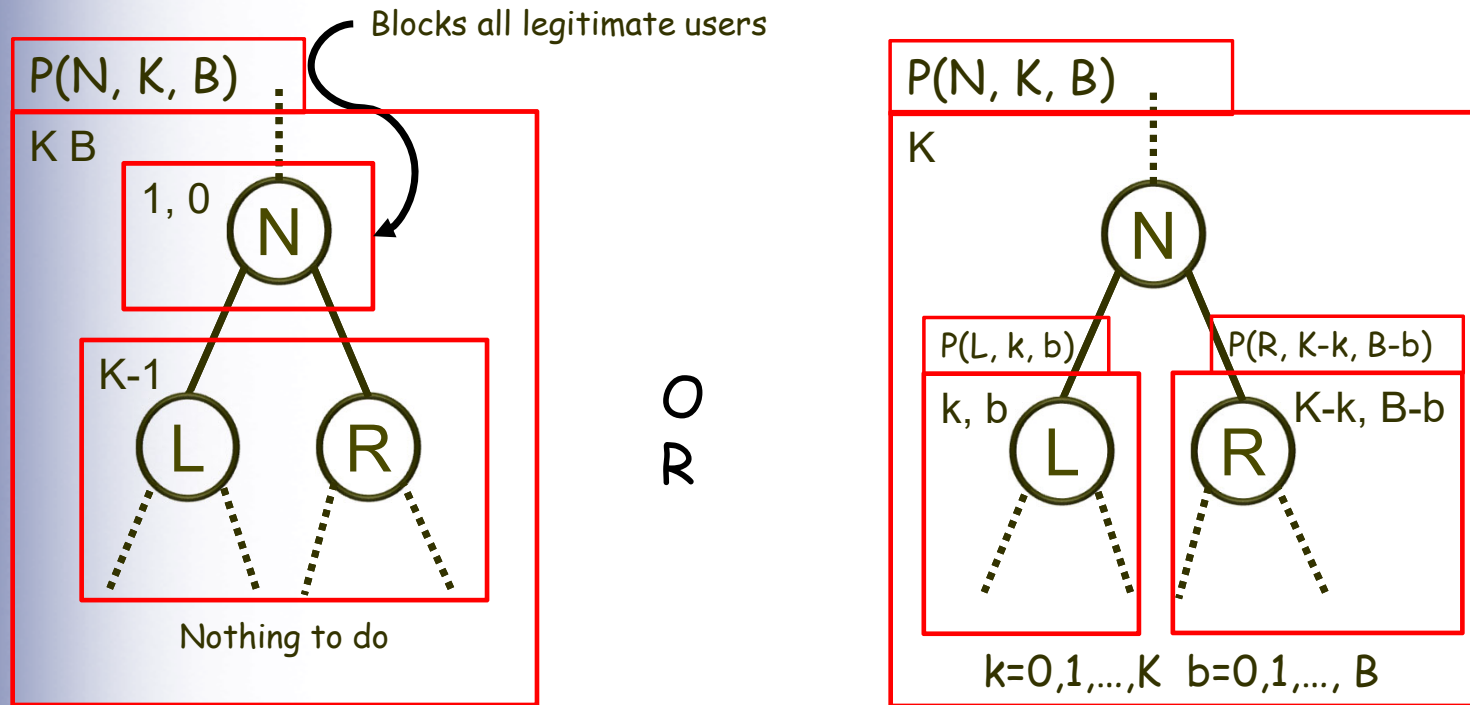
# Simplifying the Topology



- Remove nodes with no fork.



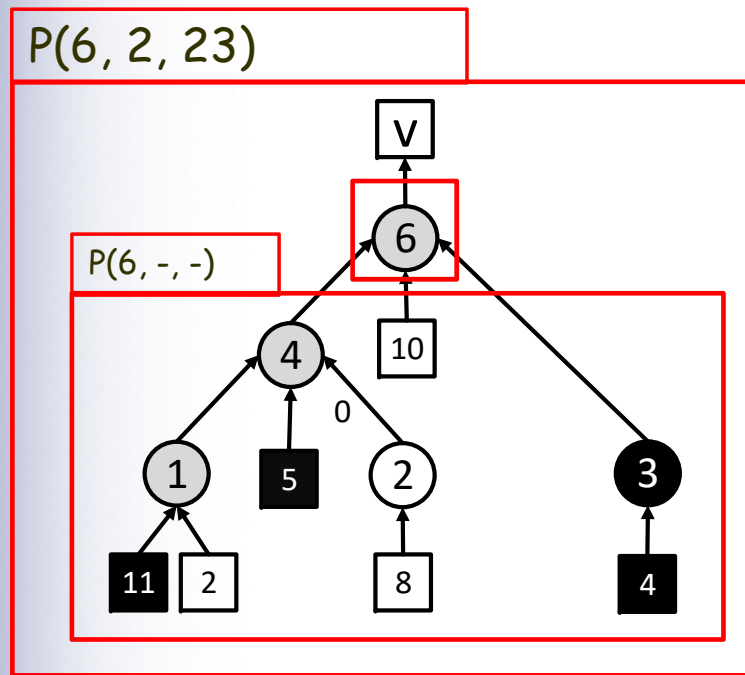
# A dynamic programming solution



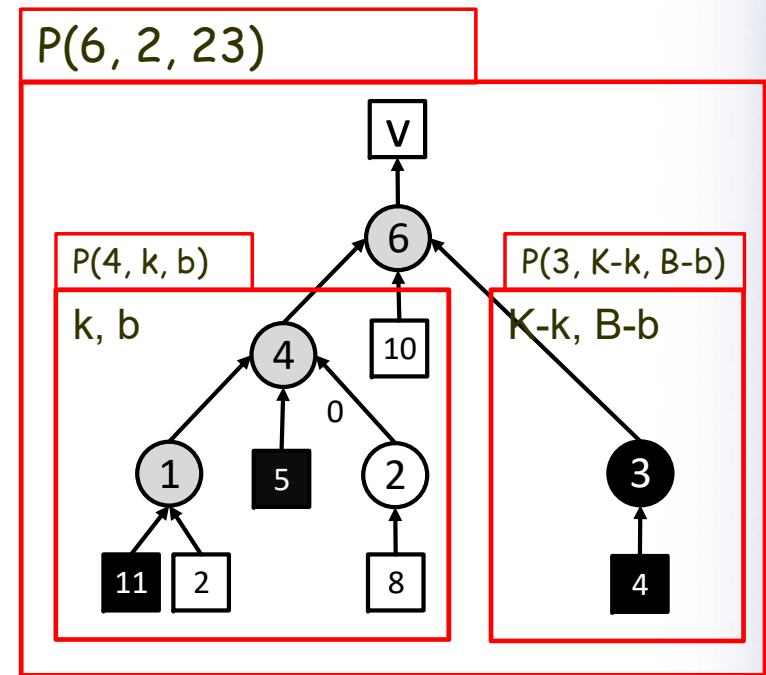
In subtree rooted by  $N$  :  $P(N, K, B) =$  Minimum blocked LU for  $K$  filters by yielding  $B$  traffic.  
 Complexity:  $O(N(KB)^{(D-1)})$



# A Dynamic Programming Solution: An Example



$$C = 20$$



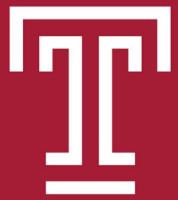
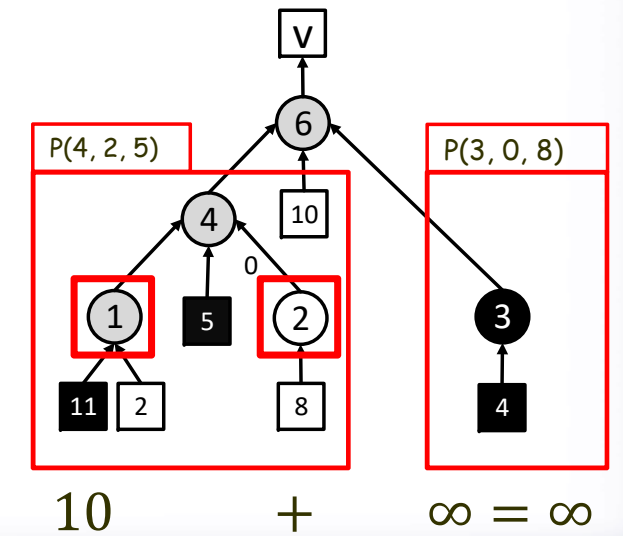
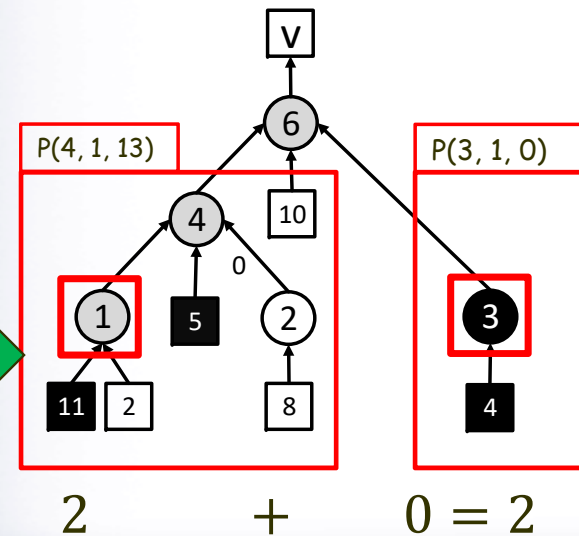
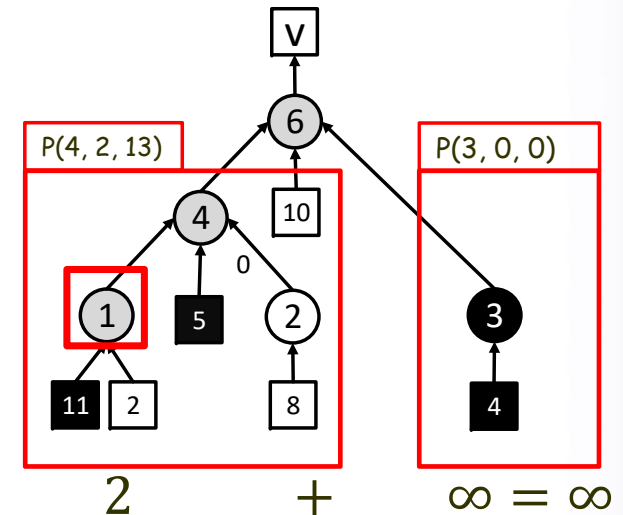
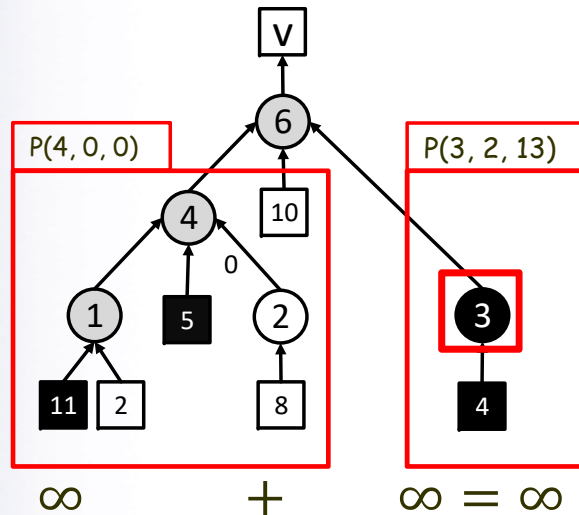
$$k=0,1,2 \quad b=0,1,2,\dots,13=23-10$$

$$C = 2$$





# A DP Solution: An Example



Minimum

# Simulation: Random Tree Generation

## Tree(d, n)

If  $d=0$

Return.

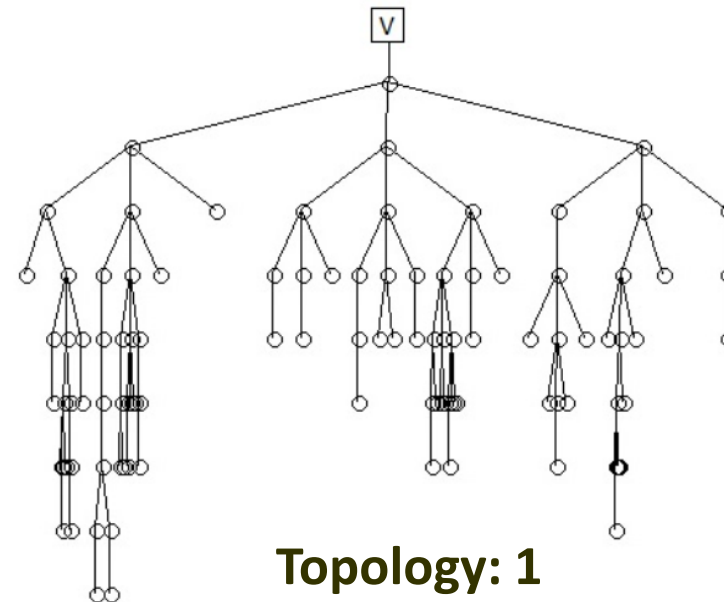
Else

For  $i=0$  to  $\text{rand}[0, \Delta]$

Create node  $c_i$ .

Make  $c_i$  child of  $n$ .

Tree( $d-1, c_i$ )



**Topology: 1**

## **Topology: 1**

# of nodes : 100

Internal user probability: 0.1

Max node degree= 3

## **Topology: 2**

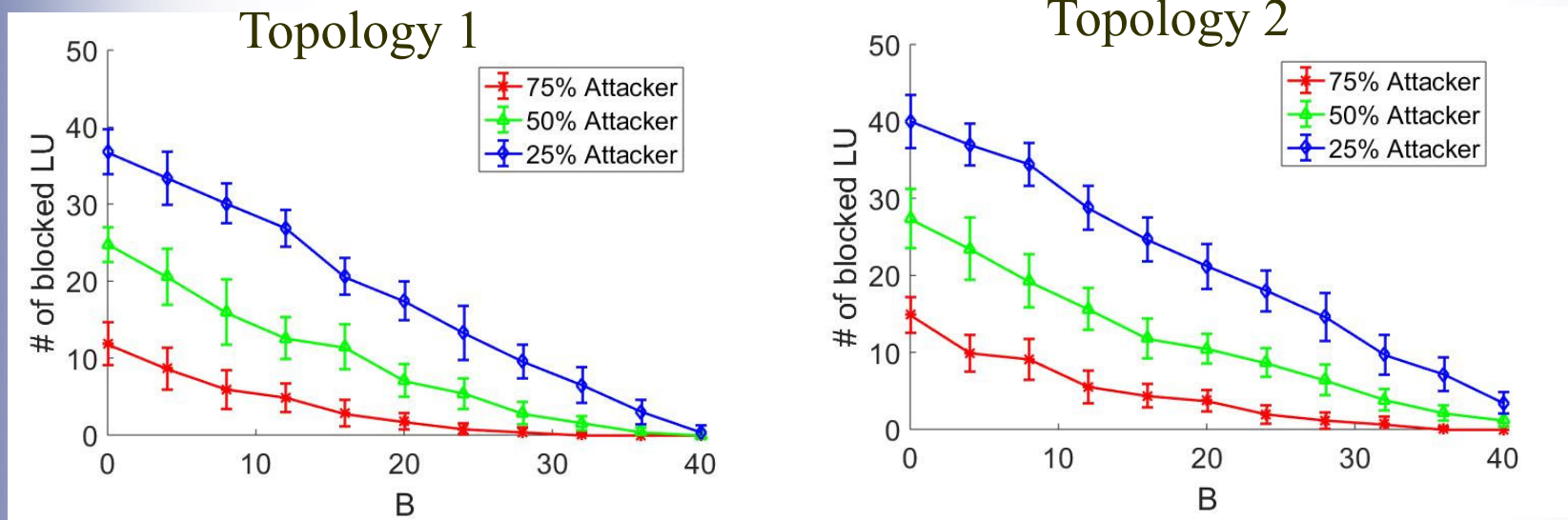
# of nodes : 400

Internal user probability: 0.1

Max node degree= 20



# Simulation: Different Number of Filters

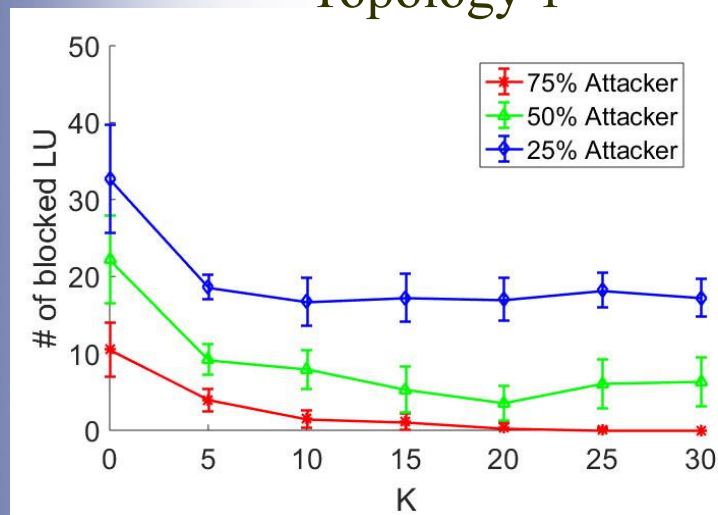


Number of blocked LUs decreases linearly with the increase of B.  
 The higher the number of attackers the higher the number of blocked LU.

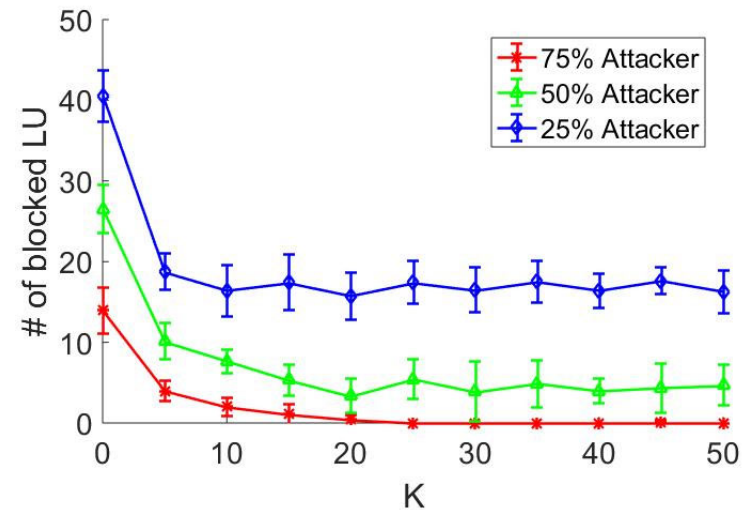


# Simulation: Different Incoming Bandwidth

## Topology 1



## Topology 2

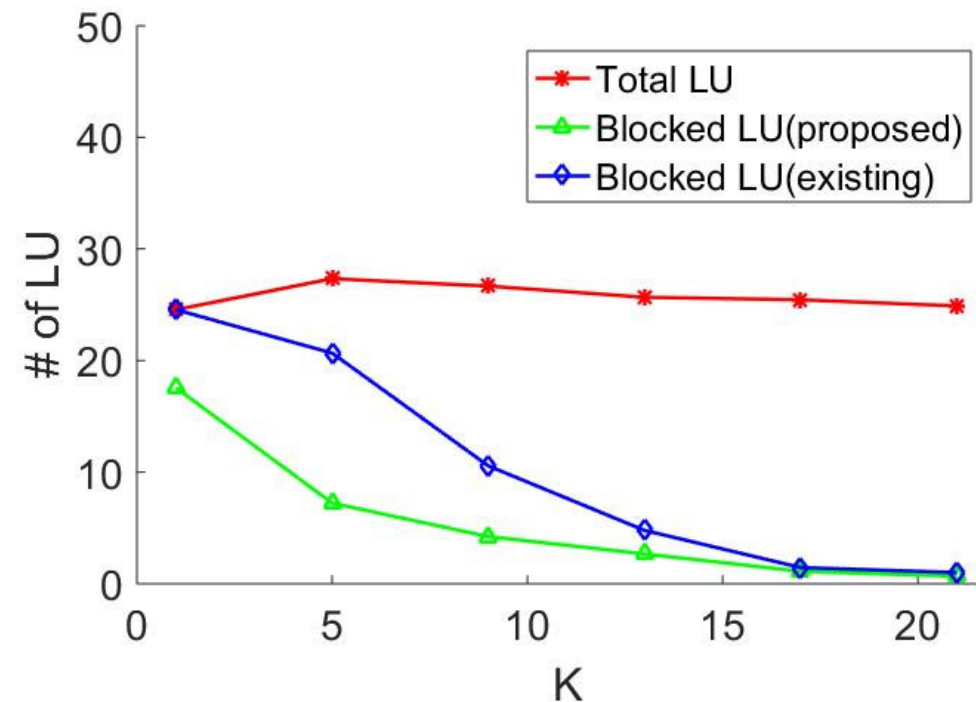


Number of blocked LUs decreases with the increase of K.  
 Number of blocked LU becomes stable after a certain value of K.





## Simulation: Compare with Existing Work

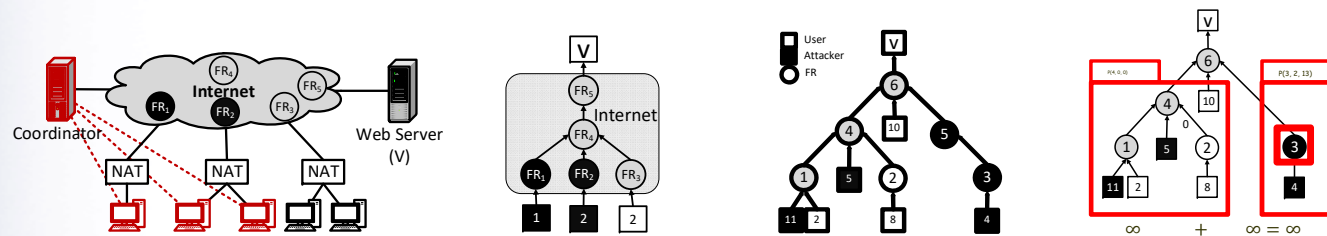


Number of blocked LUs in proposed system is less than the existing system.



# Summary

By allowing some attacker to reach victim we can significantly decrease the number of blocked legitimate user.



Thank You

Q & A !!!

