

The background features a close-up of a black wrought-iron gate with a large, gold-colored octagonal seal of Temple University. The seal depicts a classical building with columns and is surrounded by the text 'TEMPLE UNIVERSITY' and 'PHILADELPHIA'. To the left, a red banner with white Chinese characters is partially visible. A semi-transparent red banner is overlaid across the middle of the image, containing the title and authors' names in white text.

# Trajectory Scheduling for Timely Data Report in Underwater Wireless Sensor Networks

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# Agenda

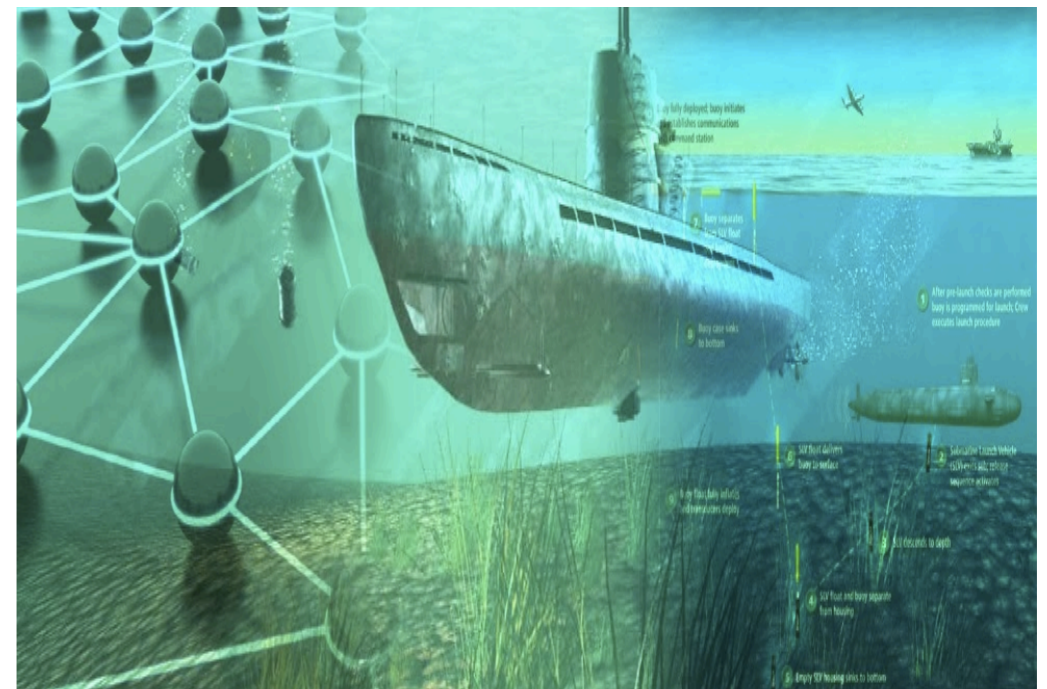
- Motivation
- Model and problem formulation
- The idea of the proposed algorithms
- Performance evaluations





# Motivation

- Ocean monitoring:
  - Research projects.
  - Pollution and disaster monitoring
    - 2004 Indian ocean earthquake
    - 2011 Japan nuclear disaster
  - Military and homeland security





# Motivation

- **Traditional method (ocean-bottom)**
  - Acoustic communication are typically expensive (US\$10k or more).
  - Can only get the data after the monitoring mission.
  - The amount of data can be recorded is limited.
- **New method (Autonomous underwater vehicles (AUV))**
  - Optical communication (cheap)
  - Collect data periodically.
  - The amount of data can collected is huge.





# Related works

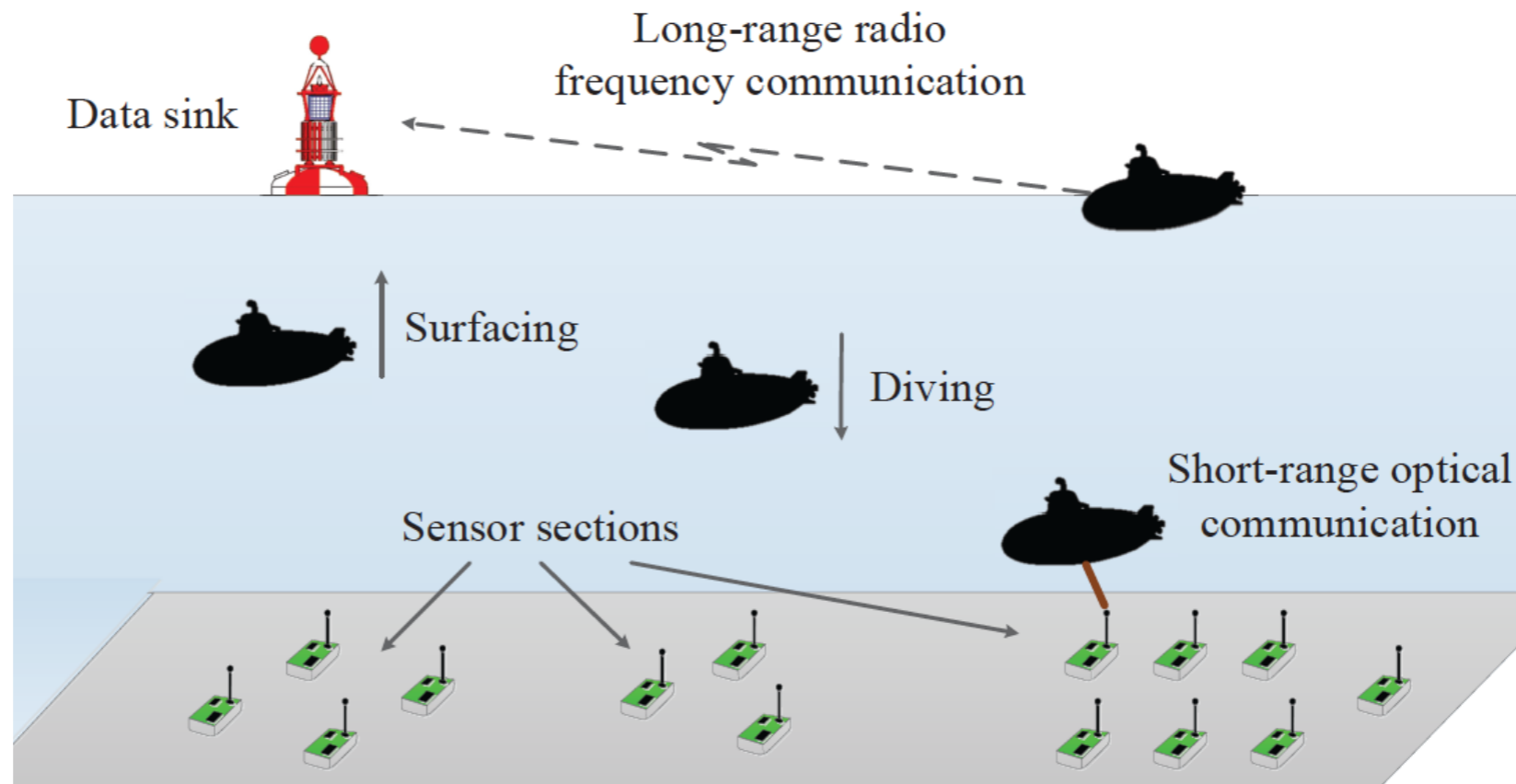
- Maximizing the value of sensed information in underwater wireless sensor networks via an autonomous underwater vehicle (INFOCOM 2014)
- Data Collection and Event Detection in the Deep Sea with Delay Minimization (SECON 2015)





# Network model and Problem

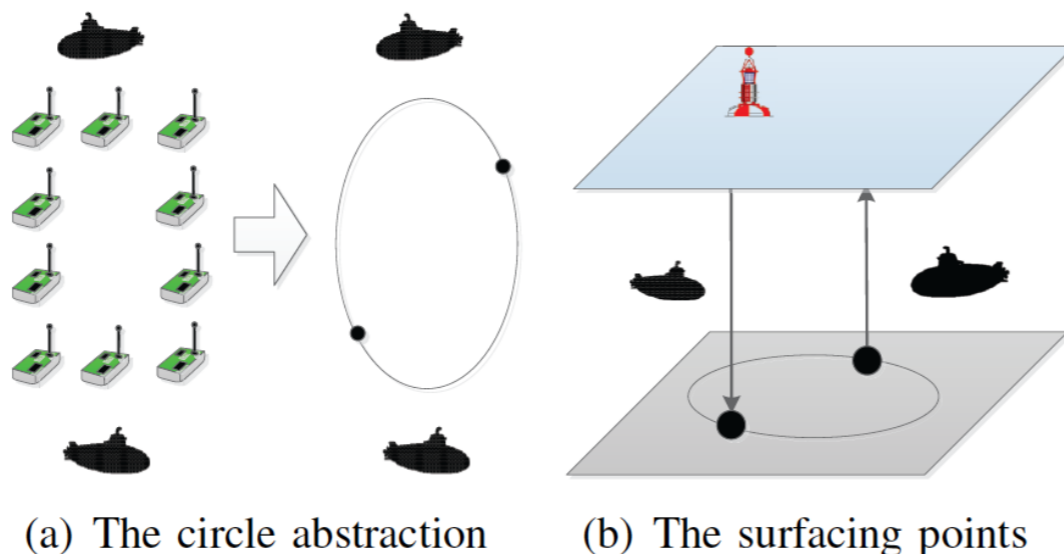
- Multiple homogeneous AUVs data collection.
  - Data are uniformly distributed with a fixed generation rate.
  - Problem: Collect all the data before their deadline.





# Network model (cont'd)

- 2D Sensor Circle Abstraction
  - Each AUV periodically collects data.



$C$  : the cycle circumference  
 $L$  : the searching space depth  
 $k$  : the number of AUVs

- If we have a larger AUV resurfacing frequency, the AUV can bring a node's data **to the water surface more quickly**.
- However, a node's data needs to **wait the next AUV for a longer time**, since resurfacings take additional time.





# Problem formulation

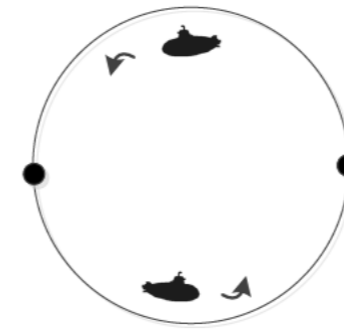
- Background:
  - Surfacing and diving is costly.
- Trajectory Scheduling problem:
  - cyclic tours with lengths  $\{c_1, c_2, \dots, c_m\}$ .
  - The number of homogeneous AUVs  $\{k_1, k_2, \dots, k_m\}$ ,
  - How to minimize **the whole amount of surfacing of AUVs**, under the constraint that **all the data** generated by the sensors can be transmitted to the sink within the deadline,  $T$ .



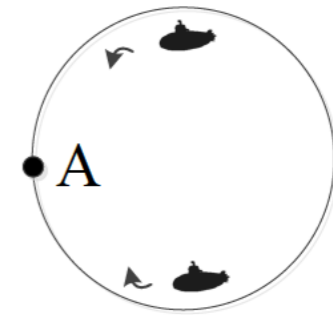


# Challenges

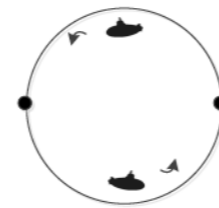
- Trajectory scheduling
  - One AUV
  - One cycle
    - Same direction
    - Different directions
- Several cycles
  - Detour or not?
  - How?



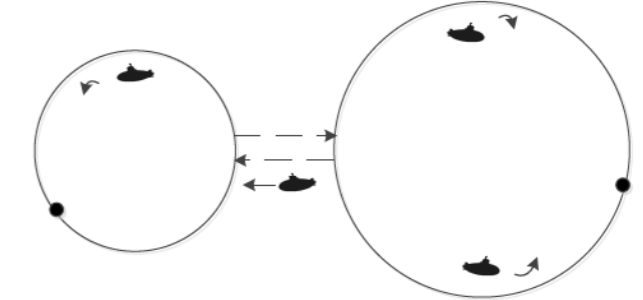
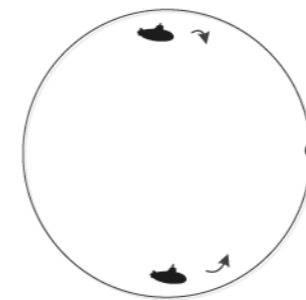
(a) Same direction



(b) Different directions



(c) Without detouring



(d) With detouring



# Within one cycle

- Same direction:
- $k$  AUVs evenly distributed in a cycle.

If we have multiple AUVs ( $k$  AUVs), then we can evenly distribute these AUVs on the cycle.

- The reporting delay is bounded by:

$$\frac{1}{v} \left( \frac{2md + c}{k} + \frac{c}{m} + d \right)$$

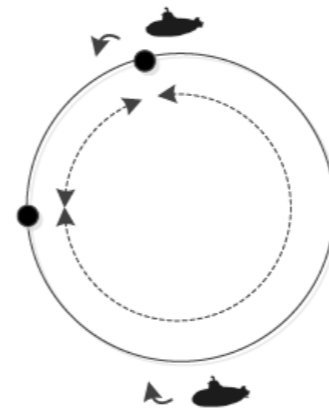
$C$  : the cycle circumference  
 $L$  : the searching space depth  
 $k$  : the number of AUVs  
 $m$  : the number of surfacing in a cycle



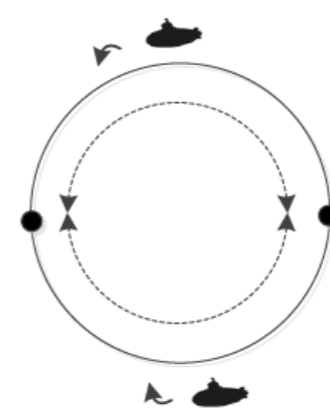


# Within one cycle

- Different directions:
- Encountered AUVs can exchange data.
- Save one surfacing



(b) Schedule 2



(c) Schedule 3

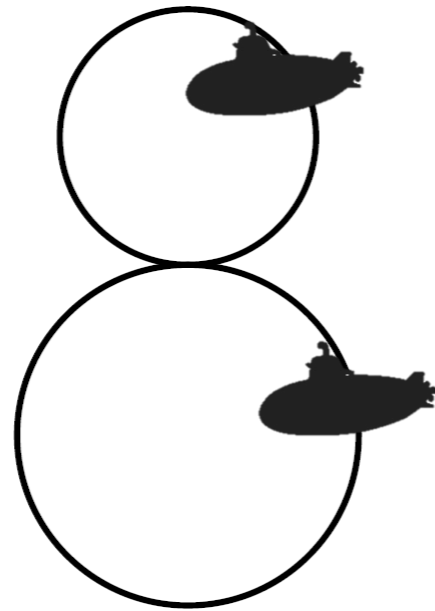
- Theorem: For a tour with an even number,  $k$ , of AUVs, the optimal schedule for minimizing the amount of surfacing is to assign  $k/2$  AUVs in one direction to surface every time of  $T - (c/k+d)/v$ .



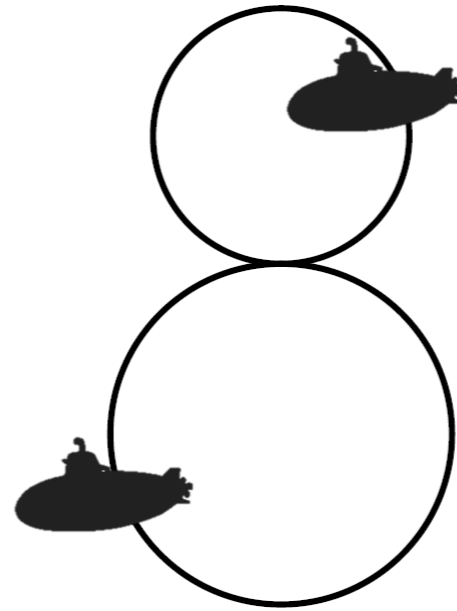


# Within several cycles

- Why do we use of **multiple small cycles** instead **only one large cycle** to collect data?



Scheduling 1



Scheduling 2

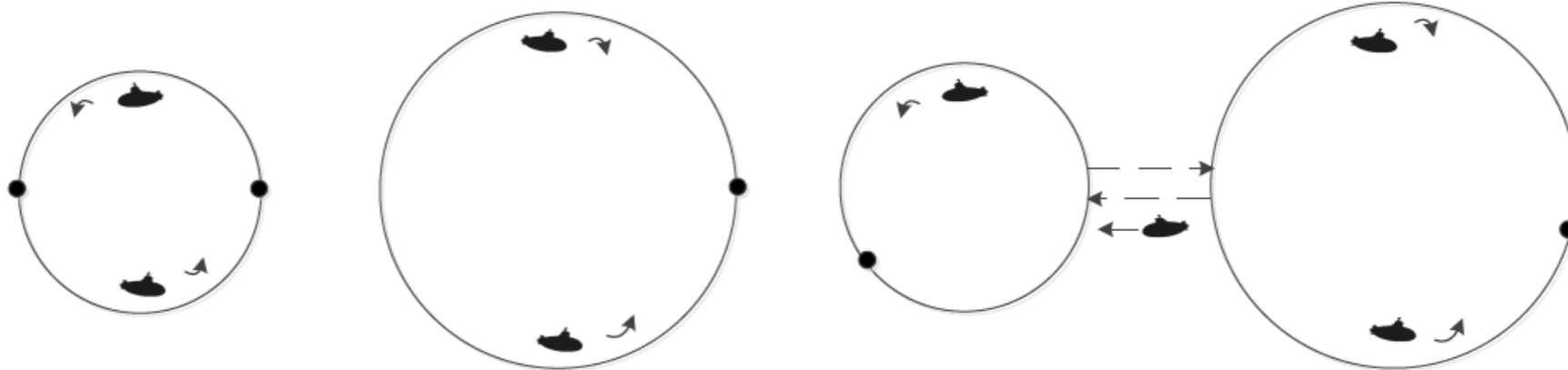
Theorem. Scheduling 2 is always no worse than Scheduling 1, due to more balanced cycling tasks among AUVs





# Within several cycles

- The real situation:
- To collaboratively schedule AUVs in two cycles will have some cost.



(c) Without detouring

(d) With detouring

Scheduling 1

Scheduling 2

There exists a trade-off of benefit and the detour distance.





# Within several cycles

- Algorithm
- 1. Calculate the cost of schedule AUVs in two AUVs individually.
- 2. Merge the two cycles into a big cycle, with some detour distance, using the previous method
- Compare the 1 and 2.

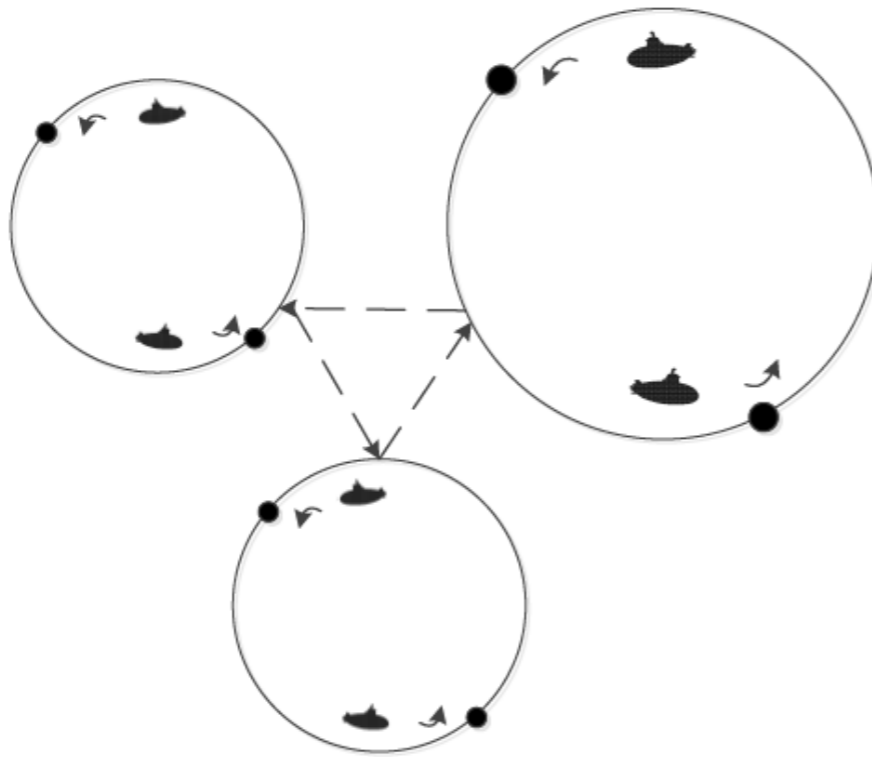
Theorem. There exists an  $1 + 2l/d$  approximation ratio between the schedule in this algorithm and the optimal solution in the merged cycle.



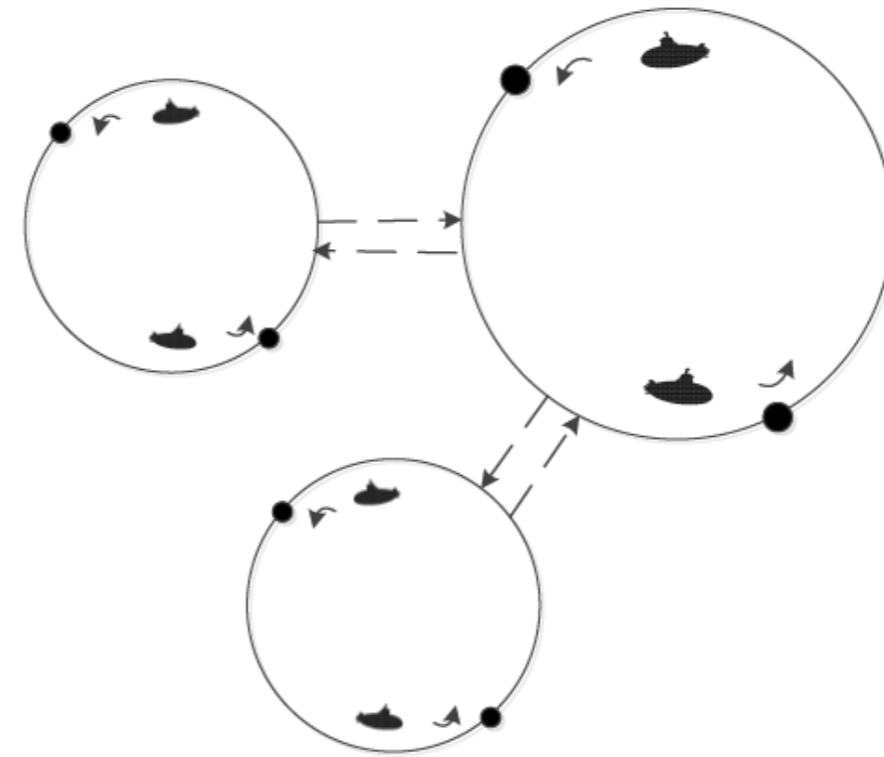


# Within several cycles

- Three cycles merge



(a) Back and forth merging



(b) Circle merging

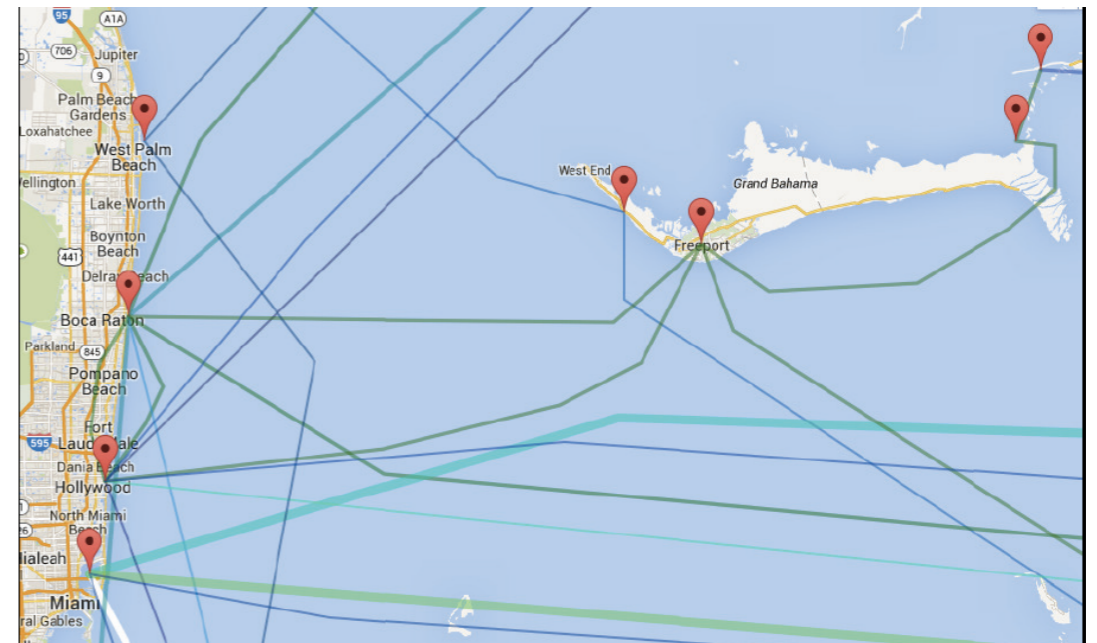




# Experiment setting

- Experiments setting
  - AUV speeds
    - 20 knots, 16 knots, 12 knots (diving, moving, surfacing)
  - 20 AUVs
  - The depth of the sea is 3682 m
  - Sensors are uniformly distributed

- The oil pipeline at Florida, USA



- BDNSi, Mid Atlantic Crossing (MAC) , GlobeNet, COLUMBUS II, III, WASACE, Americas II, cable of the Americas and BAHAMAS-2.







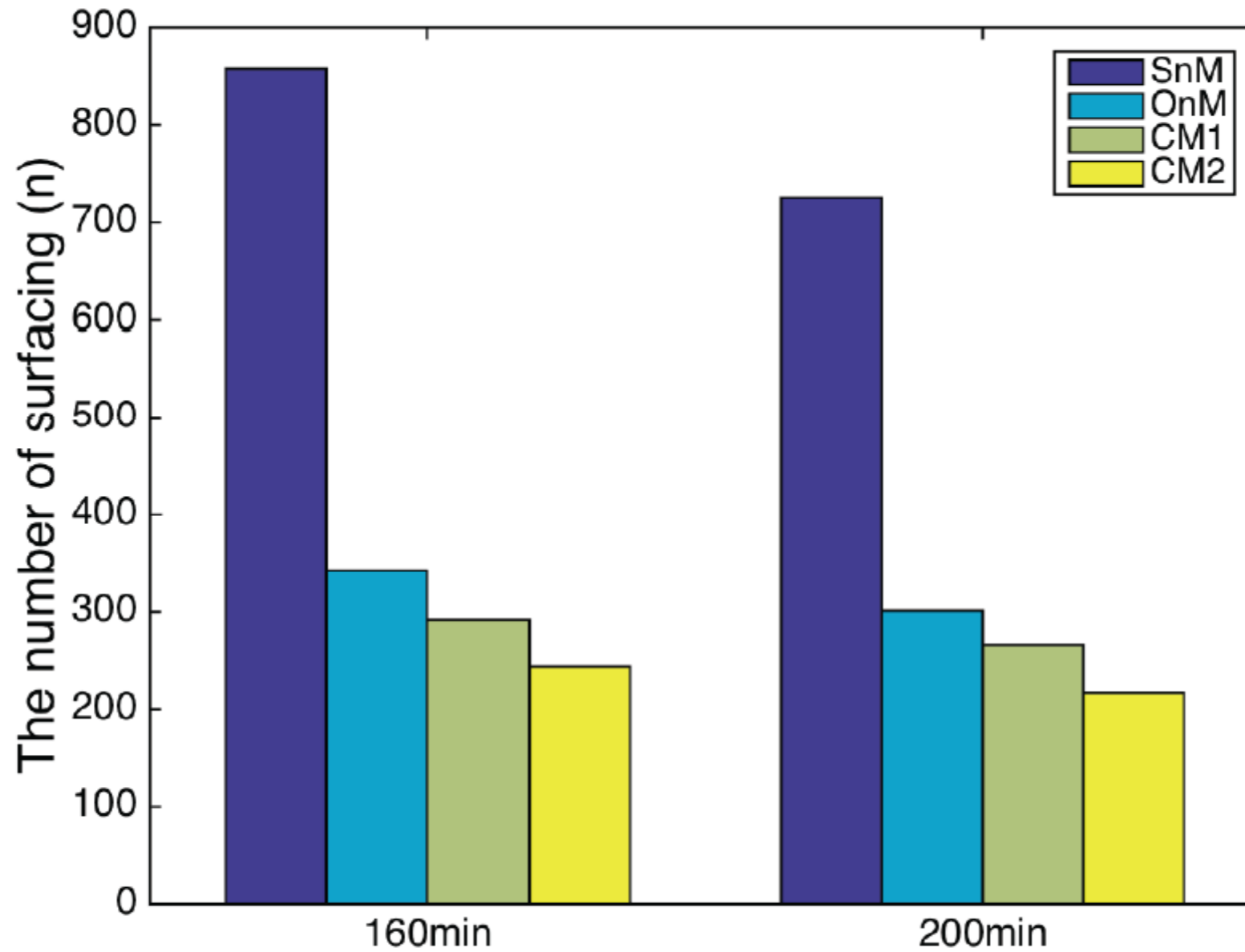
# Algorithm comparison

- Algorithms:
  - SnM: same movement direction without merging.
  - OnM: different movement directions without merging
  - CM1: only two cycle merging
  - CM2 consider the 3-cycle merging





# Experiment result





# Conclusions

- We investigate the homogeneous autonomous underwater vehicles (AUVs) trajectory schedule problem in under water sensor networks (UWSNs), considering the time constraint.
- The different scheduling methods.
  - Different moving direction within one cycle.
  - The cycle merging.





# Thank you!

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